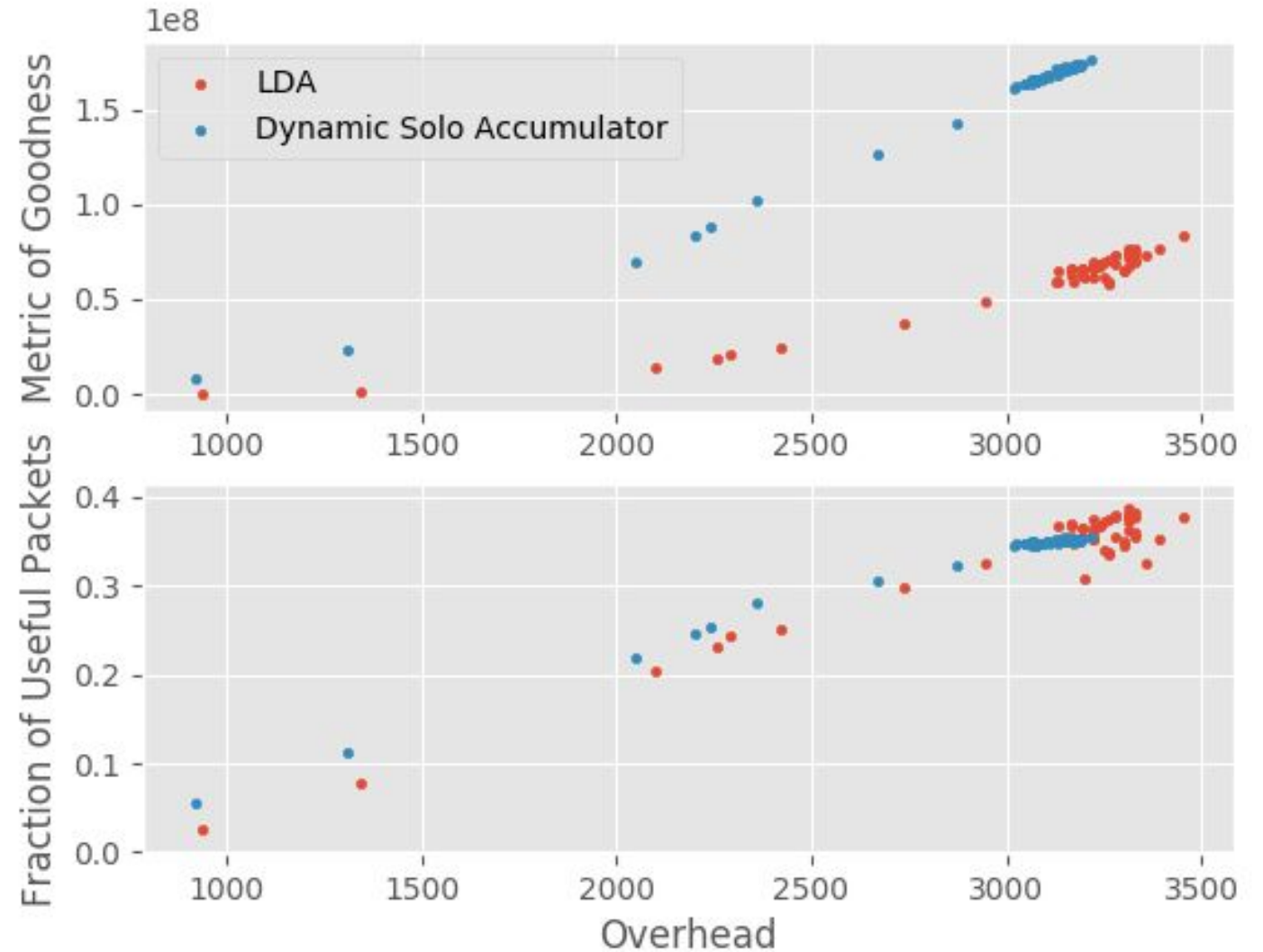


SpiNet



Window Adaptation (Single DSA)

- Time and space efficient
- Coarse grained measurement



Packet Burst/Uniform Loss

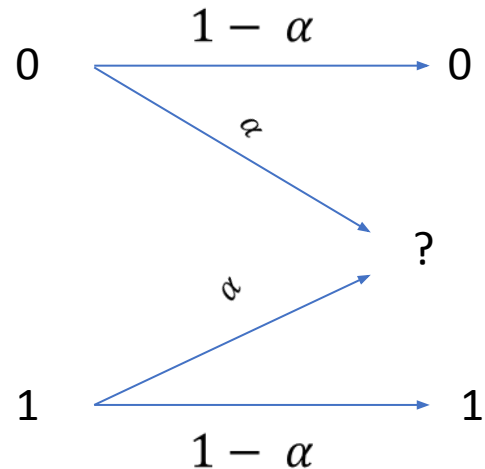
- 1. Trivial Case (single flow)

Packet Burst/Uniform Loss

- 2. Slightly looser (multiple-flow round robin case)

Packet Burst/Uniform Loss

- 3. Pure randomness case
- Modeled as a binary erasure channel
- Gödel number sequencing
- Cantor pairing
- Approximation?

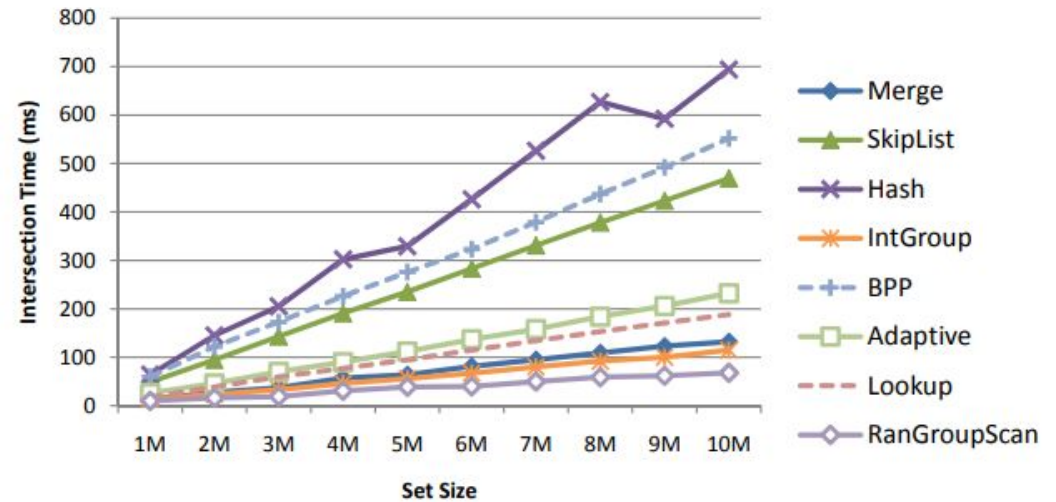


Packet Burst/Uniform Loss

- Modify outgoing packet with 'in-link' sequence number (1 Mb/s overhead). Problem reduces to [1].

Window Adaptation (Multiple DSA)

- 1. Fine grained measurement
- Header space analysis
- Locality sensitive hashing
- Accumulate counters
- 2. Fast Set intersection

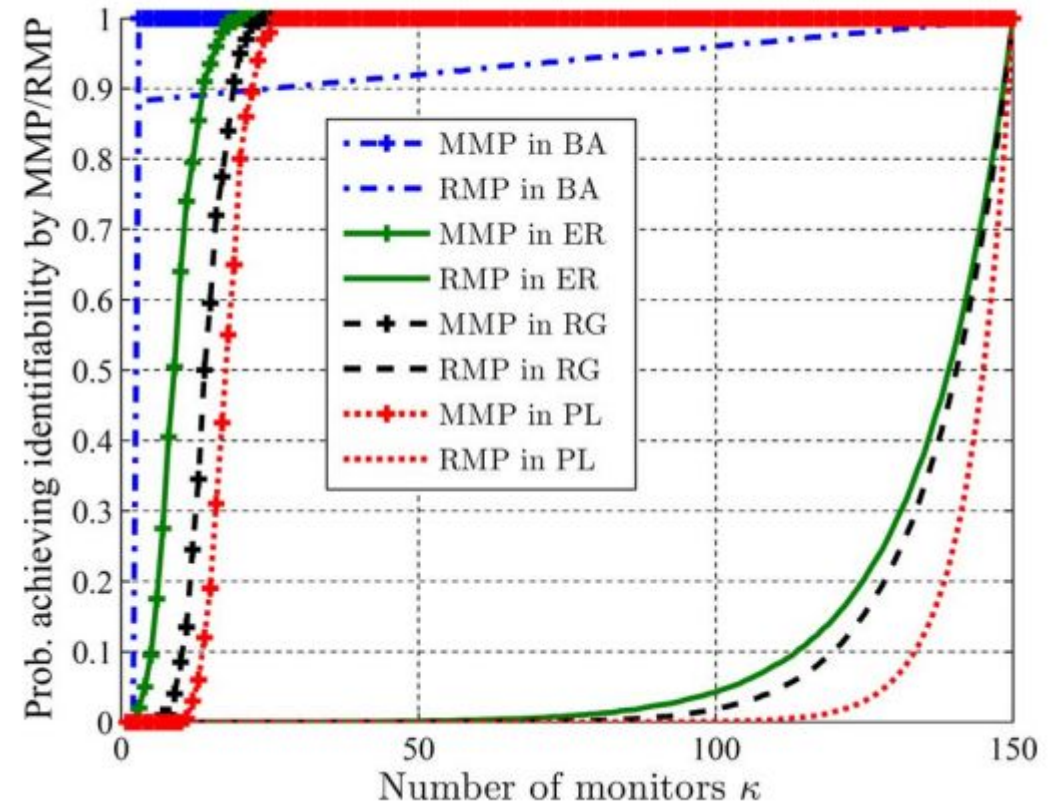


Optimal Deployment

Algorithm 1: Minimum Monitor Placement (MMP)

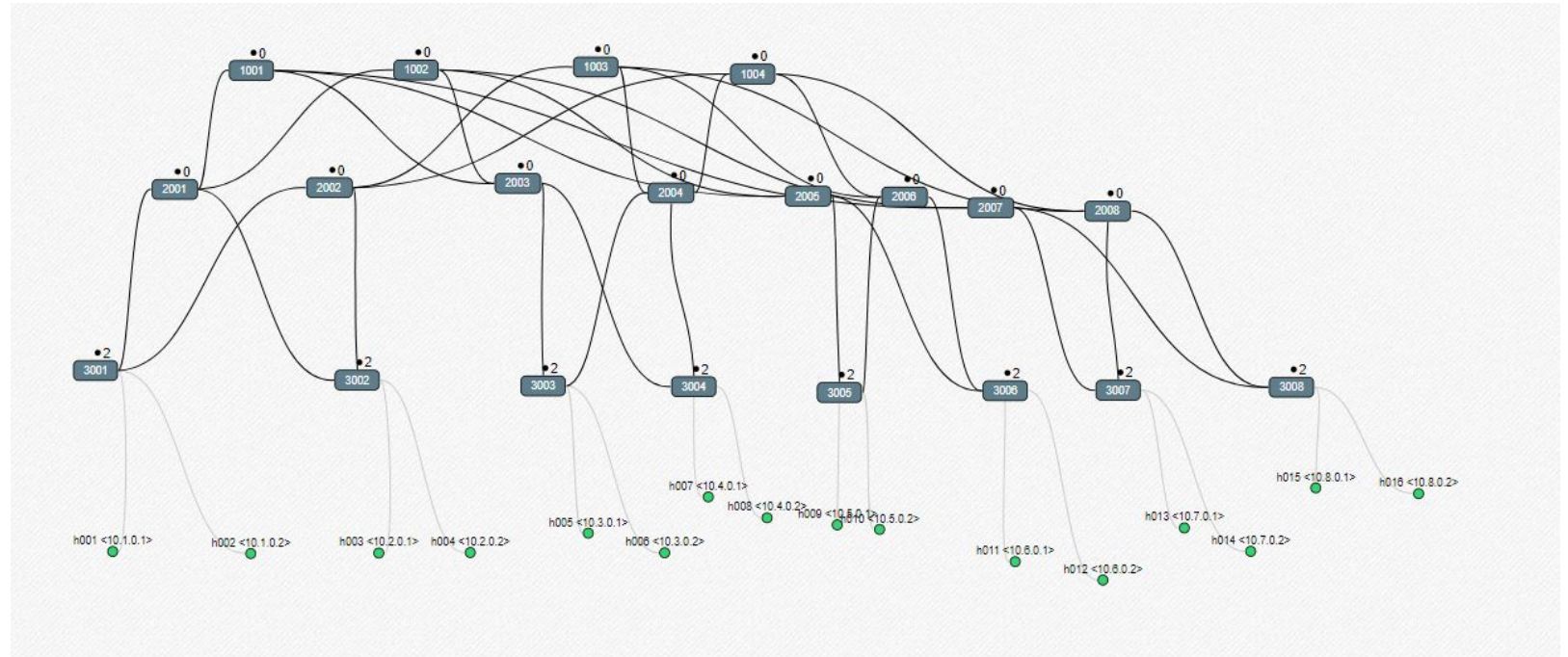
input: Connected graph \mathcal{G}
output: A subset of nodes in \mathcal{G} as monitors

- 1 choose all the nodes with degree less than 3 as monitors;
- 2 partition \mathcal{G} into biconnected components $\mathcal{B}_1, \mathcal{B}_2, \dots$;
- 3 **foreach** biconnected component \mathcal{B}_i with $|\mathcal{B}_i| \geq 3$ **do**
- 4 partition \mathcal{B}_i into triconnected components $\mathcal{T}_1, \mathcal{T}_2, \dots$;
- 5 **foreach** triconnected component \mathcal{T}_j of \mathcal{B}_i with $|\mathcal{T}_j| \geq 3$ **do**
- 6 **if** $0 < s_{\mathcal{T}_j} < 3$ **and** $s_{\mathcal{T}_j} + M_{\mathcal{T}_j} < 3$ **then**
- 7 randomly choose $3 - s_{\mathcal{T}_j} - M_{\mathcal{T}_j}$ nodes in \mathcal{T}_j that are neither separation vertices nor monitors as monitors;
- 8 **end**
- 9 **end**
- 10 **if** $0 < c_{\mathcal{B}_i} < 3$ **and** $c_{\mathcal{B}_i} + M_{\mathcal{B}_i} < 3$ **then**
- 11 randomly choose $3 - c_{\mathcal{B}_i} - M_{\mathcal{B}_i}$ nodes in \mathcal{B}_i that are neither cut-vertices nor monitors as monitors;
- 12 **end**
- 13 **end**
- 14 **if** the total number of monitors $\kappa < 3$ **then**
- 15 randomly choose $3 - \kappa$ nonmonitor nodes as monitors;
- 16 **end**



Erdős–Rényi (ER) graphs, Random Geometric (RG) graphs, Barabási–Albert (BA) graphs, and Random Power Law (PL) graphs

Evaluation



- What is the memory/bandwidth usage, and how do we compare to existing solutions (FlowRadar, LDA, NetSight, etc.)?
- What is the memory usage per switch for different deployment scenarios?
- How fast is the loss detection?
- Is the information provided useful (can we use this data to improve performance and improve rules)?
- Extra:
 - Under realistic network traffic, how effective is the inference?